Quantum Krylov algorithms for ground state energy approximation

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#### **Motivation**

Estimate ground state energy of quantum Hamiltonian.

**Abstract perspective:** approximate lowest eigenvalue of Hermitian matrix.

Goal is classically challenging due to exponential Hilbert space dimension<sup>\*</sup>.

\*Assuming general, hard case: sufficiently entangled, supported on exponentially-many basis states, etc.

#### **Example applications:**

- Quantum chemistry.
- Condensed matter physics.
- Nuclear physics.
- High-energy physics.



#### Lanczos method

= classical method for approximating lowest eigenvalues.

#### High-level idea:

- 1.
- Initial guess  $|\psi_0\rangle \Rightarrow H|\psi_0\rangle \Rightarrow \dots \Rightarrow H^{D-1}|\psi_0\rangle$ (**H**, **S**) = project *H* onto span[ $|\psi_0\rangle, H|\psi_0\rangle, H^2|\psi_0\rangle, \dots, H^{D-1}|\psi_0\rangle$ ] 2.



Lowest eigenvalue of (**H**, **S**) i.e., of  $\mathbf{H}\mathbf{v} = \lambda \mathbf{S}\mathbf{v}$ , approximates lowest eigenvalue of H 3.

Krylov space

#### Lanczos method

Caveat: typically in classical Lanczos(-like) methods, would orthogonalize along the way... challenging in quantum implementations.

Advantage: converges exponentially with D (in  $\infty$  precision arithmetic).

**Disadvantage:** classically, requires storing entire statevectors  $H^i | \psi_0 \rangle \Rightarrow$  exponential overhead.

Can we construct a quantum version that mitigates statevector overhead while keeping fast convergence?<sup>1</sup>

<sup>1</sup> Klymko *et al.*, PRX Quantum 3, 020323 (2022); Epperly *et al.*, SIAM J. Mat. An. Appl. 43, 1263-1290 (2022); and many more!

### Quantum "Lanczos method" = "Quantum Krylov"

Options for generating Krylov space: multiply  $|\psi_0\rangle$  by...

- Powers of H same as original Lanczos  $\Rightarrow$  nontrivial on quantum but possible in principle.<sup>1</sup>
- $e^{-Hk dt}$  this version claimed "Qlanczos."
- $e^{iHk dt}$  many good options e.g. Trotterization, qubitization, etc.
- $T_k(H)$  arises naturally from block encoding.

Will focus on last two in this talk.

<sup>1</sup>Seki and Yunoki, PRX Quantum 2, 010333 (2021)

#### Quantum Krylov with real time-evolutions

Majority of works have focused on Krylov states generated by real time-evolution:

$$V = [|\psi_0\rangle, \quad U|\psi_0\rangle, \quad U^2|\psi_0\rangle, \dots, U^{D-1}|\psi_0\rangle] \text{ for } U = e^{iHdt}$$

Need to estimate

 $\mathbf{H_{jk}} = \langle \psi_0 | (U^j)^{\dagger} H U^k | \psi_0 \rangle,$  $\mathbf{S_{jk}} = \langle \psi_0 | (U^j)^{\dagger} U^k | \psi_0 \rangle$ 

for each j, k = 0, 1, ..., D - 1.

# Estimating matrix elements (simple version)

Targets:  $\mathbf{H}_{\mathbf{jk}} = \langle \psi_0 | (U^j)^{\dagger} H U^k | \psi_0 \rangle$ ,  $\mathbf{S}_{\mathbf{jk}} = \langle \psi_0 | (U^j)^{\dagger} U^k | \psi_0 \rangle$ .

Can approach using Hadamard test:\*



 $\text{Yields } \langle X \rangle_a = Re\left[ \langle \psi_0 | (U^j)^{\dagger} P U^k | \psi_0 \rangle \right], \quad \langle Y \rangle_a = Im\left[ \langle \psi_0 | (U^j)^{\dagger} P U^k | \psi_0 \rangle \right]$ 

\*Cortes and Gray, 2021.

## Estimating matrix elements (better version)

Change target slightly:  $\mathbf{U}_{\mathbf{jk}} = \langle \psi_0 | (U^j)^{\dagger} U U^k | \psi_0 \rangle$ ,  $\mathbf{S}_{\mathbf{jk}} = \langle \psi_0 | (U^j)^{\dagger} U^k | \psi_0 \rangle$ .

Suppose Hamiltonian preserves particle number (Hamming weight)...

$$|0\rangle^{N} \qquad \text{prep } |\psi_{0}\rangle \qquad U^{p} \qquad (\text{prep } |\psi_{0}\rangle)^{\dagger} \qquad (\text{prep } |\psi_{0}\rangle)^{\dagger} \qquad (\text{prep } \frac{1}{\sqrt{2}}(|0\rangle^{N} + |\psi_{0}\rangle))^{\dagger} \qquad (\text{prep } \frac{1}{\sqrt{2}}(|0\rangle^{N} + |\psi|^{\dagger}))^{\dagger} \qquad (\text{prep } \frac{1}{\sqrt{2}}(|0\rangle^{N} +$$

## Quantum Krylov with real time-evolutions

#### Summary:

- Estimate  $H_{ik}$ ,  $S_{ik}$  via Hadamard(-ish) tests and repeated sampling.
- Depending on Hamiltonian, can sometimes avoid controlled time-evolutions using symmetries.<sup>1</sup>
- Advantage: can use crude approximations for time-evolution to get low circuit depth.
- Disadvantage: time-evolution always approximated more accuracy requires more depth.

<sup>1</sup>Cortes and Gray, Phys. Rev. A 105, 022417 (2022).

The most accurate simulations of time-evolution require Hamiltonian input as block encoding.<sup>1</sup>

**Block encoding:** for *H* on *n* qubits (s.t.  $||H|| \le 1$ ), find *U* on m + n qubits s.t.



$$U^{2} = 1 \qquad (RU)^{j} = \left( \begin{array}{c} T_{j}(H) & \cdot \\ \cdot & \cdot \end{array} \right)$$

Brief proof:  
1. Let 
$$H | \lambda \rangle = \lambda | \lambda \rangle$$
 and  $U = | G \rangle \langle G | \otimes H + ...$   
2.  $\Rightarrow U | G \rangle | \lambda \rangle = \lambda | G \rangle | \lambda \rangle + \sqrt{1 - \lambda^2} | \perp \rangle \Rightarrow U \sim \begin{pmatrix} \lambda & \cdot \\ \sqrt{1 - \lambda^2} & \cdot \end{pmatrix}$   
3.  $U^2 = 1 \Rightarrow U \sim \begin{pmatrix} \lambda & \sqrt{1 - \lambda^2} \\ \sqrt{1 - \lambda^2} & -\lambda \end{pmatrix} \Rightarrow RU \sim \begin{pmatrix} \lambda & \sqrt{1 - \lambda^2} \\ -\sqrt{1 - \lambda^2} & \lambda \end{pmatrix}$   
reflection rotation

$$(RU)^{j} = \left( \begin{array}{cc} T_{j}(H) & \cdot \\ \cdot & \cdot \end{array} \right)$$

 $\Rightarrow$  Can use block encoding to exactly construct  $T_j(H) | \psi_0 \rangle$ 

**Recall:** Lanczos method ~ project *H* onto

$$span\{ |\psi_0\rangle, H |\psi_0\rangle, H^2 |\psi_0\rangle, \dots, H^{D-1} |\psi_0\rangle \}$$
  
= span{  $|\psi_0\rangle, T_1(H) |\psi_0\rangle, T_2(H) |\psi_0\rangle, \dots, T_{D-1}(H) |\psi_0\rangle$ 

 $\mathbf{H}_{\mathbf{jk}} = \langle \psi_0 \,|\, T_j(H) H T_k(H) \,|\, \psi_0 \rangle$ 

$$= \frac{1}{4} \left( \left\langle T_{i+j+1}(H) \right\rangle_{0} + \left\langle T_{|i+j-1|}(H) \right\rangle_{0} + \left\langle T_{|i-j+1|}(H) \right\rangle_{0} + \left\langle T_{|i-j-1|}(H) \right\rangle_{0} \right) \right)$$

$$\mathbf{S_{jk}} = \langle \psi_0 | T_j(H) T_k(H) | \psi_0 \rangle = \frac{1}{2} \left( \left\langle T_{i+j}(H) \right\rangle_0 + \left\langle T_{|i-j|}(H) \right\rangle_0 \right)$$

for j, k = 0, 1, 2, ..., D - 1. In other words, need to estimate

 $\langle T_k(H) \rangle_0 = \langle \psi_0 | T_k(H) | \psi_0 \rangle$ 

for k = 0, 1, 2, ..., 2D - 1.

Since 
$$(RU)^{j} = \begin{pmatrix} T_{j}(H) & \cdot \\ \cdot & \cdot \end{pmatrix}$$

 $\Rightarrow \quad \langle T_k(H) \rangle_0 = \langle \psi_0 | T_k(H) | \psi_0 \rangle = (\langle G | \otimes \langle \psi_0 | )(RU)^k (| G \rangle \otimes | \psi_0 \rangle)$ 

$$= \begin{cases} (\langle G | \otimes \langle \psi_0 | \rangle (UR)^{\lfloor k/2 \rfloor} R (RU)^{\lfloor k/2 \rfloor} (|G\rangle \otimes |\psi_0\rangle) & \text{if } k \text{ is even,} \\ \underbrace{(\langle G | \otimes \langle \psi_0 | \rangle (UR)^{\lfloor k/2 \rfloor}}_{\langle \Psi_k |} U \underbrace{(RU)^{\lfloor k/2 \rfloor} (|G\rangle \otimes |\psi_0\rangle)}_{|\Psi_k\rangle} & \text{if } k \text{ is odd .} \end{cases}$$

That yields the necessary circuits: for each k = 0, 1, 2, ..., D - 1...



# **Regularizing the Krylov space**

- Either real-time or block-encoding Krylov methods yield noisy estimates of  $(\mathbf{H}, \mathbf{S})$ .
- Want to solve  $\mathbf{H}\mathbf{v} = \lambda \mathbf{S}\mathbf{v}$ .
- Ill-conditioned if S is ill-conditioned  $\Rightarrow$  need to regularize.
- "Canonical orthogonalization" or "thresholding": project (H, S) onto eigenspaces of S above threshold  $\epsilon$ .



• S is metric in Krylov space  $\Rightarrow$  choose threshold  $\sim$  noise rate  $\Rightarrow$  discards vectors compatible with 0.

# **Error analysis**

- Real-time: analyzed in Epperly et al., SIAM J. Mat. An. Appl. 43, 1263-1290 (2022).
- Block-encoding: in our paper, modification of Epperly's analysis.
- Three main error terms:

error from Krylov space + error from thresholding + error from noise



- First two terms high-level idea of proof:
  - Krylov space = span[ $|\psi_0\rangle$ ,  $T_1(H) |\psi_0\rangle$ ,  $T_2(H) |\psi_0\rangle$ , ...,  $T_{D-1}(H) |\psi_0\rangle$ ] = poly<sub>D-1</sub>(H)  $|\psi_0\rangle$
  - $\Rightarrow$  Best poly approx to delta function at  $E_0$  = approx ground space projector in Krylov space.
  - Thresholding  $\Rightarrow$  perturbation of Chebyshev expansion coefficients of projector.

## **Error analysis**

"In practice" results: 1 to reach energy error  $\mathscr{C}$ , require...

Krylov space dimension 
$$D = \Theta \left[ \left( \log \frac{1}{|\gamma_0|} + \log \frac{1}{\mathscr{C}} \right) \min \left( \frac{1}{\mathscr{C}}, \frac{1}{\Delta} \right) \right],$$

Measurements per dimension 
$$M = \Theta\left(\frac{1}{\mathscr{E}^2} + \frac{1}{\mathscr{E}|\gamma_0|^4}\right)$$

where  $\gamma_0$  = initial state overlap with low-energy subspace,  $\Delta$  = spectral gap.

<sup>1</sup> "In practice" because theoretical bound only guarantees first term in M is  $\mathscr{C}^{-p}$  for  $p \in [2,3]$ ; p = 2 is based on numerics.

<sup>2</sup> Above is for block-encoding; real-time Krylov space analysis is similar.

## Comparing theory to numerics



 $\delta$  free, best choice is  $\delta = \Theta(\max(\text{target error}, \Delta))$ 





Questions?

Will Kirby, Mario Motta, and Antonio Mezzacapo, Quantum 7, 1018 (2023),

https://quantum-journal.org/papers/q-2023-05-23-1018/.